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WHAT IS CLAIMED IS:

1. In a decoder used for error detection of a codeword, a method for evaluating a single error location polynomial coefficient generated from said codeword in a cell corresponding to said single error location polynomial coefficient, said method comprising the acts of:

receiving an error location polynomial coefficient corresponding to said codeword;

multiplying said error location polynomial coefficient, on a first clock cycle corresponding to the processing of said codeword, by a Galois field multiplier having a negative exponent, wherein said negative exponent is a function of a stage number (j) corresponding to said cell and the length of said codeword (N), said act of multiplying resulting in a cell output; and

iteratively multiplying said cell output, for a subsequent N minus one clock cycles, by a Galois field multiplier having a positive exponent, wherein said positive exponent is a function of said stage number (j).

- 2. The method of claim 1 in which said cell is a Chien search cell of a Chien search block.
- 3. The method of claim 1 in which said cell is a Forney algorithm cell of a Forney algorithm block.
- 4. The method of claim 1 in which said cell is a Forney algorithm cell of a Chien/Forney block.
 - 5. The method of claim 1 in which said decoder is a Reed-Solomon decoder.
- 30 6. In a decoder used for error detection of a codeword, a method for evaluating a single error location polynomial coefficient generated from said codeword in a cell corresponding to said single error location polynomial

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coefficient, the operations of said method controlled by the parameters of the equations:

$$\begin{aligned} \mathbf{X}_i &= \sum_{j=0}^t & \boldsymbol{\varLambda}_j \; \boldsymbol{\alpha}^{-j(N-1)} & \text{for i = 0} \\ \\ \mathbf{X}_i &= \sum_{j=0}^t & \boldsymbol{\varLambda}_j \; \boldsymbol{\alpha}^j & \text{for i = 1, 2, ..., (N-1)} \end{aligned}$$

5 in which,

 X_i is the i- th error location root,

 $\it t$ is one less than the total number of coefficients corresponding to an error location polynomial,

 $ec{\Lambda}_{j}$ is the j-th error polynomial coefficient,

N is the codeword length,

 α is a Galois field element, and

j is a stage number corresponding to said single error location polynomial coefficient

7. An apparatus for evaluating a single error location polynomial coefficient generated from a codeword in a cell corresponding to said single error location polynomial coefficient, said apparatus incorporated within a decoder used for error detection of said codeword, said apparatus comprising:

means for receiving an error location polynomial coefficient corresponding to said codeword;

means for multiplying said error location polynomial coefficient, on a first clock cycle corresponding to the processing of said codeword, by a Galois field multiplier having a negative exponent, wherein said negative exponent is a function of a stage number (j) corresponding to said cell and the length of said codeword (N), said act of multiplying resulting in a cell output; and

means for iteratively multiplying said cell output, for a subsequent N minus one clock cycles, by a Galois field multiplier having a positive exponent, wherein said positive exponent is a function of said stage number (j).

- 8. The apparatus of claim 7 in which said cell is a Chien search cell of a Chien search block.
- 9. The apparatus of claim 7 in which said cell is a Forney algorithm cell of a Forney algorithm block.
 - 10. The apparatus of claim 7 in which said cell is a Forney algorithm cell of a Chien/Forney block.
- 10 11. The apparatus of claim 7 in which said decoder is a Reed-Solomon decoder.